

On η -Expansion in NbE and Type Casts

Andreas Abel
 Department of Computer Science
 Ludwig-Maximilians-University Munich
 andreas.abel@ifi.lmu.de

13 October 2011

This small note justifies the asymmetry in the definition of the η -expansion functions \uparrow and \downarrow for function types in the context of normalization by evaluation for dependent types [ACD07, Abe10, ACP11].

$$\begin{aligned} (\uparrow^{\text{Fun } A F} n) a &= \uparrow^{F a} n (\downarrow^A a) \\ (\downarrow^{\text{Fun } A F} f) x &= \downarrow^{F \uparrow^A x} f (\uparrow^A x) \end{aligned}$$

The asymmetry $F a$ vs. $F \uparrow^A x$ can be derived from a view of \uparrow and \downarrow as embedding-projection pair, or up- and downcast functions, with a symmetric definition:

$$\begin{aligned} (\uparrow^{\text{Fun } A' F'} f) a' &= \uparrow_{F (\downarrow_{A'} a)}^{F' a'} f (\downarrow_{A'} a') \\ (\downarrow^{\text{Fun } A' F'} f') a &= \downarrow_{F a}^{F' (\uparrow_{A'} a)} f' (\uparrow_{A'} a) \end{aligned}$$

Upcast $\uparrow_A^{A'}$ casts a value of type A to one of type A' ; downcast $\downarrow_A^{A'}$ casts in the opposite direction, from A' to A . The definition is symmetric, we have for all A, A' ,

$$\uparrow_A^{A'} = \downarrow_{A'}^A$$

The upcast $\uparrow^A n$ in the context of NbE casts a neutral $n \in \text{Ne}$ into A , and the downcast $\downarrow^A a$ casts a value $a \in A$ into the set of normal forms Nf . Since

$$\text{Ne} \subseteq \text{Nf} \rightarrow \text{Ne} \text{ and } \text{Ne} \rightarrow \text{Nf} \subseteq \text{Nf}$$

we can identify Ne with $\text{Fun Nf } \lambda.\text{Ne}$ and Nf with $\text{Fun Ne } \lambda.\text{Nf}$. More precisely, if we want to convert from Ne it is sufficient to convert from $\text{Nf} \rightarrow \text{Ne}$. And if we want to convert into Nf it is sufficient to convert into $\text{Ne} \rightarrow \text{Nf}$. This yields

$$\begin{aligned} (\uparrow_{\text{Ne}}^{\text{Fun } A' F'} f) a' &= (\uparrow_{\text{Fun Nf } \lambda.\text{Ne}}^{\text{Fun } A' F'} f) a' = \uparrow_{\text{Ne}}^{F' a'} f (\downarrow_{\text{Nf}}^{A'} a') \\ (\downarrow_{\text{Nf}}^{\text{Fun } A' F'} f') a &= (\downarrow_{\text{Fun Ne } \lambda.\text{Nf}}^{\text{Fun } A' F'} f') a = \downarrow_{\text{Nf}}^{F' (\uparrow_{\text{Ne}}^{A'} a)} f' (\uparrow_{\text{Ne}}^{A'} a) \end{aligned}$$

from which we can drop the redundant indices Ne and Nf to arrive at the NbE-definition of \uparrow and \downarrow .

References

- [Abe10] Andreas Abel. Towards Normalization by Evaluation for the $\beta\eta$ -Calculus of Constructions. In *FLOPS'10*, volume 6009 of *LNCS*, pages 224–239. Springer, 2010.
- [ACD07] Andreas Abel, Thierry Coquand, and Peter Dybjer. Normalization by evaluation for Martin-Löf Type Theory with typed equality judgements. In *LICS'07*, pages 3–12. IEEE CS Press, 2007.
- [ACP11] Andreas Abel, Thierry Coquand, and Miguel Pagano. A modular type-checking algorithm for type theory with singleton types and proof irrelevance. *LMCS*, 7(2:4):1–57, 2011.